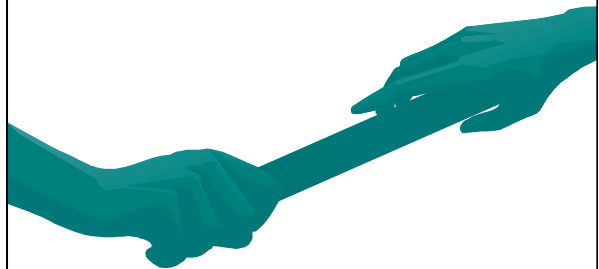


# Automata and Languages

Prof. Mohamed Hamada

Software Engineering Lab.  
The University of Aizu  
Japan

## FINITE STATE MACHINES (AUTOMATA)



### Switch Example 1

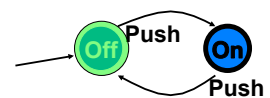


Think about the On/Off button

### Switch Example 1



The corresponding Automaton



Input:

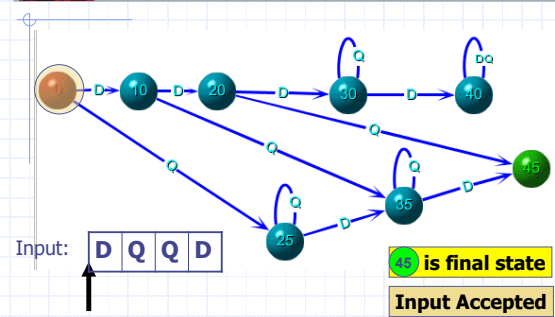
Push Push Push Push

## Vending Machine Example 2

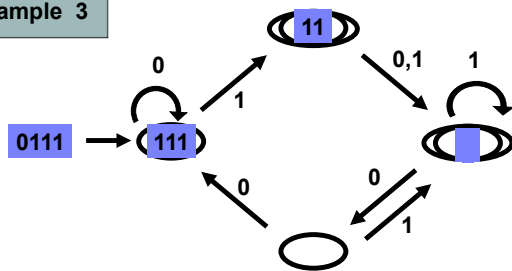


✦ Vending machine dispenses Cola for \$0.45

## Vending Machine Example 2

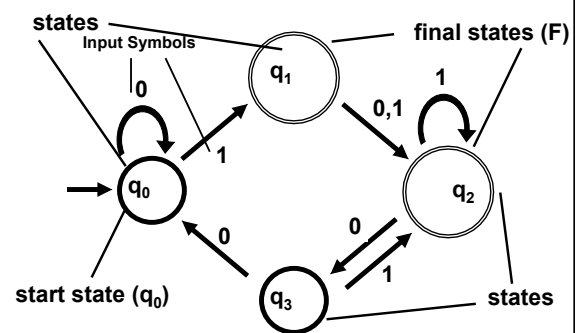


## Example 3



The machine accepts a string if the process ends in a *final state*

## Example 3



### Definitions

An **alphabet**  $\Sigma$  is a finite set of symbols  
(in Ex3,  $\Sigma = \{0,1\}$ )

A **string** over  $\Sigma$  is a finite sequence of elements of  $\Sigma$  (e.g. 0111)

For a string  $s$ ,  $|s|$  is the **length** of  $s$

The unique string of length 0 will be denoted by  $\lambda$  and will be called the **empty string**

The **reversal** of a string  $u$  is denoted by  $u^R$ .  
Example:  $(\text{banana})^R = \text{ananab}$

### Definitions

The **concatenation** of two strings is the string resulting from putting them together from left to right. Given strings  $u$  and  $v$ , denote the concatenation by  $u.v$ , or just  $uv$ .

**Example:**

jap . an = japan, QQ . DD = QQDD

Q1: What's the Java equivalent of concatenation?

The + operator on strings

Q2: Find a formula for  $|u.v|$ ?

$|u.v| = |u| + |v|$

### Definitions

If  $\Sigma$  is an alphabet,  
 $\Sigma^*$  denotes the set of all strings over  $\Sigma$ .

A **language** over  $\Sigma$  is a subset of  $\Sigma^*$

i.e. a set of strings *each* consisting of sequences of symbols in  $\Sigma$ .

### Examples

**Example1:** in our vending machine we have

$\Sigma = \{D, Q\}$

$\Sigma^* = \{\lambda,$

D, Q,  
DD, DQ, QD, QQ,  
DDD, DDQ, DQD, DQQ, QDD, QDQ, QQD, QQQ,  
DDDD, DDDQ, ... }

$L = \{u \in \Sigma^* \mid u \text{ successfully vends}\}$



**Example2:** in our switch example we have

$\Sigma = \{\text{Push}\}$

$\Sigma^* = \{\lambda,$

Push,  
Push Push,  
Push Push Push,  
Push Push Push Push, ... }

$L = \{\text{Push}^n \mid n \text{ is odd}\}$



## Definitions

A finite automaton is a 5-tuple  $M = (Q, \Sigma, \delta, q_0, F)$

$Q$  is the set of states

$\Sigma$  is the alphabet

$\delta$  is the transition function

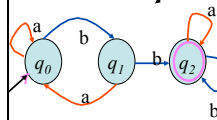
$q_0 \in Q$  is the start state

$F \subseteq Q$  is the set of final states

$L(M)$  = the language of machine  $M$   
= set of all strings machine  $M$  accepts

## Definitions

State Diagram and Table



$Q = \{q_0, q_1, q_2\}$   
 $\Sigma = \{a, b\}$   
 $F = \{q_2\}$

$\delta$	a	b
$q_0$	$q_0$	$q_1$
$q_1$	$q_0$	$q_2$
$q_2$	$q_2$	$q_2$

## FINITE STATE MACHINES (AUTOMATA)

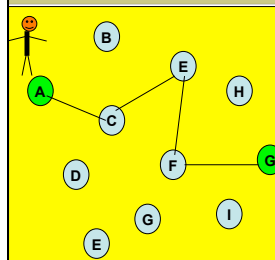
Deterministic Finite Automata (DFA)

Non-Deterministic Finite Automata with empty move ( $\lambda$ -NFA)

Non-Deterministic Finite Automata (NFA)

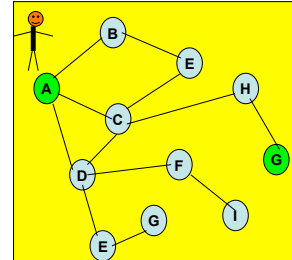
## Deterministic & Nondeterministic

Deterministic

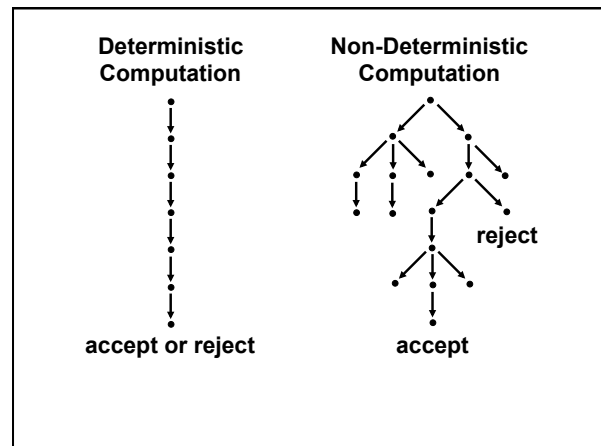
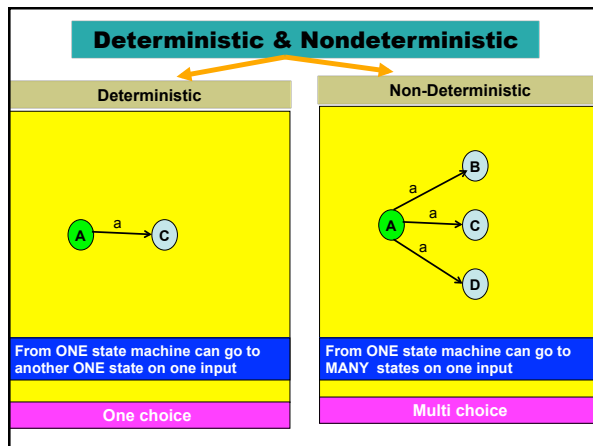


One choice

Non-Deterministic



Multi choice  $\rightarrow$  Backtrack

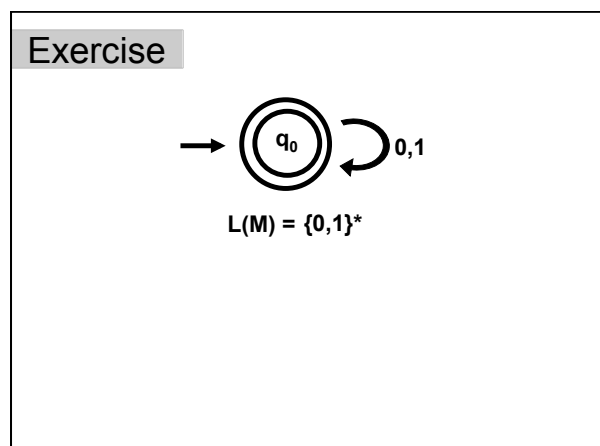
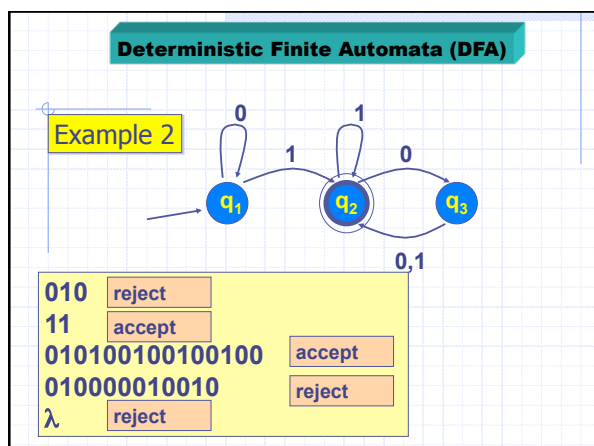
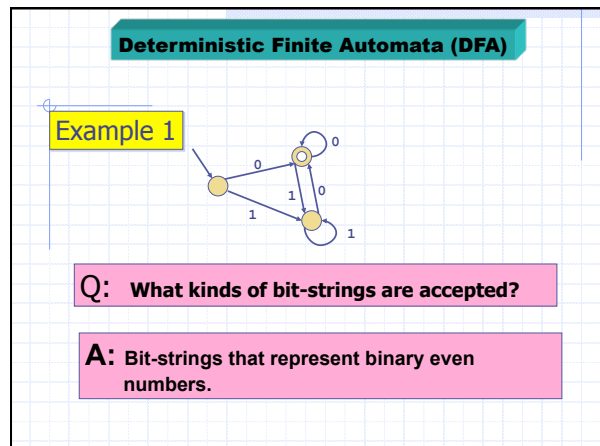
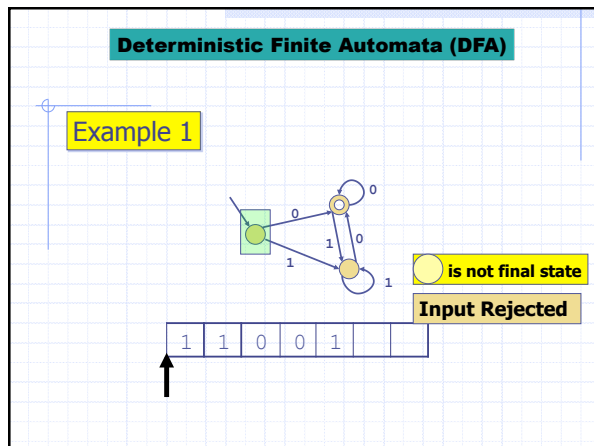


**Definitions**

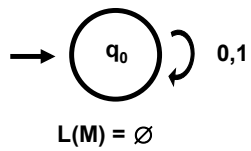
A DFA is a 5-tuple  $M = (Q, \Sigma, \delta, q_0, F)$

- $Q$  is the set of states
- $\Sigma$  is the alphabet
- $\delta : Q \times \Sigma \rightarrow Q$  is the transition function
- $q_0 \in Q$  is the start state
- $F \subseteq Q$  is the set of accept states

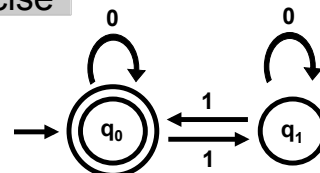
$L(M)$  = the language of machine  $M$   
 = set of all strings machine  $M$  accepts



### Exercise



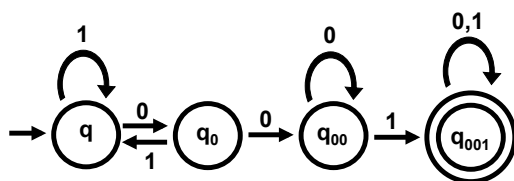
### Exercise



$L(M) = \{ w \mid w \text{ has an even number of 1s} \}$

### Exercise

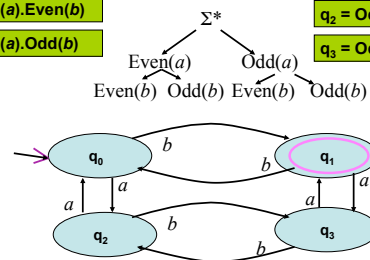
Build an automaton that accepts all and only those strings that contain 001



### Exercise

Strings over  $\{a,b\}$  containing even number of  $a$ 's and odd number of  $b$ 's.

$q_0 = \text{Even}(a).\text{Even}(b)$        $q_2 = \text{Odd}(a).\text{Even}(b)$   
 $q_1 = \text{Even}(a).\text{Odd}(b)$        $q_3 = \text{Odd}(a).\text{Odd}(b)$



## Exercise

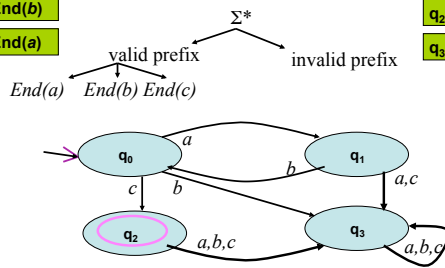
Strings over  $\{a,b,c\}$  that has the form  $(ab)^*c$

$q_0 = \text{End}(b)$

$q_1 = \text{End}(a)$

$q_2 = \text{End}(c)$

$q_3 = \text{Error}$



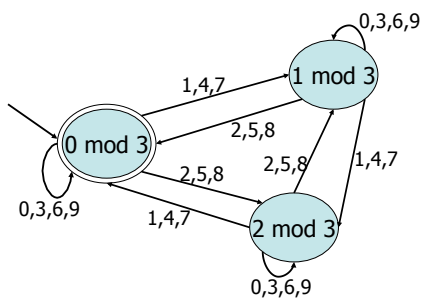
## Exercise

Design with a friend a machine that tells us when a *base-10* number is divisible by 3.

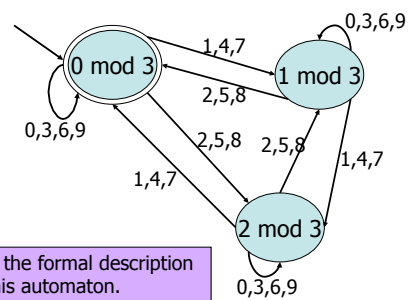
What should your alphabet be?

How can you tell when a number is divisible by 3?

## Answer



## Exercise



Find the formal description of this automaton.



## Answer

$Q = \{ 0 \bmod 3, 1 \bmod 3, 2 \bmod 3 \}$  (rename:  $\{q_0, q_1, q_2\}$ )  
 $\Sigma = \{ 0, 1, 2, 3, 4, 5, 6, 7, 8, 9 \}$   
 $q_0 = 0 \bmod 3$   
 $F = \{ 0 \bmod 3 \}$

$\delta : Q \times \Sigma \rightarrow Q$

$\delta(q_0, 2) = q_2, \delta(q_0, 9) = q_0, \delta(q_1, 2) = q_0,$   
 $\delta(q_1, 7) = q_2, \delta(q_2, 3) = q_2, \delta(q_2, 5) = q_1.$

Question :  $\delta(q_i, j) = ?$

$$\delta(q_i, j) = q_{(i+j) \bmod 3}$$